An Introduction to Formal Real Analysis, Rutgers University, Fall 2025, Math 311H

Lecture 4: NonConvergence

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1 A Great Debate

In the 19th century, when people were still confused about limits in general, there was a specific sequence that they fought over:

```
1, -1, 1, -1, 1, -1, ...
```

Does this converge or not???

Some people said, well, you go up one then down one, so on average, it's around zero. So the limit should be 0.

Other people said: look at the partial sums of the sequence:

$$1 = 1$$

$$1 + (-1) = 0$$

$$1 + (-1) + 1 = 1$$

$$1 + (-1) + 1 + (-1) = 0$$

Ah so the partial sums alternate between 1 and 0, so maybe the limit is 1/2!

This is the kind of trouble you get into if you don't have a community consensus on what the definitions of things are, namely, if you don't have a rigorous definition of limit. Luckily, we now do!

With that definition, this sequence does not have a limit ("diverges"). In order to express that formally, let's make a new definition:

```
def SeqConv (a : \mathbb{N} \rightarrow \mathbb{R}): Prop := \existsL, SeqLim a L
```

That is, a sequence a "converges" (without specifying to what), if there indeed exists some real number L so that a converges to that number.

In this level, then, our goal is to prove:

```
Goal : ¬SeqConv a,
```

given the assumption $ha : \forall n$, $a n = (-1 : \mathbb{R})^n$. (Note that the "not" symbol, \neg , is obtained by typing $\$

How do you prove the negation of something? Logically speaking, you would say, well, if that thing did happen, then we'd have a contradiction. Technically, $\neg P$ is definitionally equivalent to: $P \longrightarrow False$.

So we could start our proof by reminding ourselves of this fact, by typing: change SeqConv a \rightarrow False

That will change the Goal to: SeqConv a \rightarrow False.

In general, I would recommend the following protocol:

- 1) Think. Do scratchwork, get a vague sense of how to piece the argument together.
- 2) Then prove it formally. You're NOT done! Just because you proved it formally doesn't *necessarily* mean that you really understand what's going on! The goal of mathematics is *not* that it just works, but rather that you **understand** exactly what's going on. So step 3 is:
- 3) Give a natural language proof that explains (to yourself, as much as to anyone else, including me) what's going on and why.

Using only knee jerk reactions (if you see $P \to Q$ in the Goal, write intro. If you see \exists in a hypothesis, write choose. If you see a definition in the Goal or a hypothesis, write change), we got as far as:

```
Objects: a : \mathbb{N} \to \mathbb{R}, L : \mathbb{R} Assumptions: ha : \forall(n : \mathbb{N}), a n = (-1)^ n, hL : \forall \varepsilon > 0, \exists \mathbb{N}, \foralln \geq \mathbb{N}, |a n - L| < \varepsilon Goal: False
```

Remember the Engineer and Mechanic: the Engineer gets to specify any tolerance, and the Mechanic has to guarantee that we'll be within that tolerance "eventually", that is, for all large enough ${\tt n}$, as measured by the lower bound ${\tt N}$.

Key Idea: Find an ε that is too tight a tolerance, and the Mechanic will never be able to get within that specification. Because the sequence alternates between 1 and -1, which differ by 2, the tolerance $\varepsilon = 1$ would already be enough, but just to be safe, let's give ourselves a little more room, and set $\varepsilon = 1/2$.

Note: Had we set $\varepsilon = 1.5$, that would *not* work. We can't rule out the possibility that L= 0, and both values 1 and -1 *are* indeed within that tolerance. So we won't find a contradiction that way.

After a bunch of computation, we reached the following: |1 - L| < 1/2 and |-1 - L| < 1/2 Let's try to come up with a contradiction from that. |1 - L| < 1/2 Let's try to come up with a contradiction from that |1 - L| < 1/2 Let's try to come up with a contradiction from that. |1 - L| < 1/2 Let's try to come up with a contradiction from that. |1 - L| < 1/2 Let's try to come up with a contradiction from that. |1 - L| < 1/2 Let's try to come up with a contradiction from that. |1 - L| < 1/2 Let's try to come up with a contradiction from that. |1 - L| < 1/2 Let's try to come up with a contradiction from that. |1 - L| < 1/2 Let's try to come up with a contradiction from that. |1 - L| < 1/2 Let's try to come up with a contradiction from that. |1 - L| < 1/2 Let's try to come up with a contradiction from that. |1 - L| < 1/2 Let's try to come up with a contradiction from that. |1 - L| < 1/2 Let's try to come up with a contradiction from that. |1 - L| < 1/2 Let's try to come up with a contradiction from that. |1 - L| < 1/2 Let's try to come up with a contradiction from that. |1 - L| < 1/2 Let's try to come up with a contradiction from that. |1 - L| < 1/2 Let's try to come up with a contradiction from that.

1.1 New Tools You'll Need

- Negation: If P is some Prop, then ¬P is definitionally equivalent to P → False. So you can write change P → False, either at the Goal, or at a hypothesis.
- The triangle inequality: To add the fact that $|x + y| \le |x| + |y|$ to our list of hypotheses, invoke the abs_add theorem:

```
have factName : |x + y| \le |x| + |y| := by apply abs_add
```

• Negation inside an absolute value: You may also find useful the theorem abs_neg, which can be called via:

```
have factName : |-x| = |x| := by apply abs_neg
```

Warning! Make sure the pattern |-Something| is on the left hand side. If Lean doesn't see an absolute value followed by a minus sign, abs_neg won't work!

2 Step 2: Lean solution

```
Statement (a : \mathbb{N} \to \mathbb{R}) (ha : \forall n, a n = (-1) ^ n) : \neg SeqConv a := by change SeqConv a \to False intro h change \exists L, SeqLim a L at h
```

```
choose L hL using h
change \forall \varepsilon > 0, \exists N, \forall n \geq N, |a n - L| < \varepsilon at hL
specialize hL (1/2)
have whatever : (0 : \mathbb{R}) < 1/2 := by norm_num
specialize hL whatever
choose N hN using hL
have h2N : N < 2 * N := by bound
have h2Np1 : N \le 2 * N + 1 := by bound
have f1 : |a(2 * N) - L| < 1 / 2 := by apply hN (2 * N) - L | < 1 / 2 := by apply hN (2 * N) - L | < 1 / 2 := by apply hN (2 * N) - L | < 1 / 2 := by apply hN (2 * N) - L | < 1 / 2 := by apply hN (2 * N) - L | < 1 / 2 := by apply hN (2 * N) - L | < 1 / 2 := by apply hN (2 * N) - L | < 1 / 2 := by apply hN (2 * N) - L | < 1 / 2 := by apply hN (2 * N) - L | < 1 / 2 := by apply hN (2 * N) - L | < 1 / 2 := by apply hN (2 * N) - L | < 1 / 2 := by apply hN (2 * N) - L | < 1 / 2 := by apply hN (2 * N) - L | < 1 / 2 := by apply hN (2 * N) - L | < 1 / 2 := by apply hN (2 * N) - L | < 1 / 2 := by apply hN (2 * N) - L | < 1 / 2 := by apply hN (2 * N) - L | < 1 / 2 := by apply hN (2 * N) - L | < 1 / 2 := by apply hN (2 * N) - L | < 1 / 2 := by apply hN (2 * N) - L | < 1 / 2 := by apply hN (2 * N) - L | < 1 / 2 := by apply hN (2 * N) - L | < 1 / 2 := by apply hN (2 * N) - L | < 1 / 2 := by apply hN (2 * N) - L | < 1 / 2 := by apply hN (2 * N) - L | < 1 / 2 := by apply hN (2 * N) - L | < 1 / 2 := by apply hN (2 * N) - L | < 1 / 2 := by apply hN (2 * N) - L | < 1 / 2 := by apply hN (2 * N) - L | < 1 / 2 := by apply hN (2 * N) - L | < 1 / 2 := by apply hN (2 * N) - L | < 1 / 2 := by apply hN (2 * N) - L | < 1 / 2 := by apply hN (2 * N) - L | < 1 / 2 := by apply hN (2 * N) - L | < 1 / 2 := by apply hN (2 * N) - L | < 1 / 2 := by apply hN (2 * N) - L | < 1 / 2 := by apply hN (2 * N) - L | < 1 / 2 := by apply hN (2 * N) - L | < 1 / 2 := by apply hN (2 * N) - L | < 1 / 2 := by apply hN (2 * N) - L | < 1 / 2 := by apply hN (2 * N) - L | < 1 / 2 := by apply hN (2 * N) - L | < 1 / 2 := by apply hN (2 * N) - L | < 1 / 2 := by apply hN (2 * N) - L | < 1 / 2 := by apply hN (2 * N) - L | < 1 / 2 := by apply hN (2 * N) - L | < 1 / 2 := by apply hN (2 * N) - L | < 1 / 2 := by apply hN (2 * N) - L | < 1 / 2 := by apply hN (2 * N) - L | < 1 / 2 := by apply hN (2 * N) - L | < 1 / 2 := by apply hN (2 * N) - L | < 1 / 2 := by apply hN (2 * N) - L | < 1 / 2 := by apply hN (2 * N) - L | < 1 / 2 := by apply hN (2 * N) - L | < 1 / 2 := by apply h
        N) h2N
have f2 : |a(2 * N + 1) - L| < 1 / 2 := by apply hN
         (2 * N + 1) h2Np1
have f3 : a (2 * N) = (-1) ^ (2 * N) := by apply ha (2)
have f4 : a (2 * N + 1) = (-1) ^ (2 * N + 1) := by
         apply ha (2 * N + 1)
rewrite [f3] at f1
rewrite [f4] at f2
have f5 : (-1 : \mathbb{R}) ^ (2 * \mathbb{N}) = 1 := by bound
rewrite [f5] at f1
have f6 : (-1 : \mathbb{R}) ^ (2 * \mathbb{N} + 1) = -1 := by bound
rewrite [f6] at f2
have f7 : (2 : \mathbb{R}) = |2| := by norm_num
have f8 : |(2 : \mathbb{R})| = |1 - (-1)| := by ring_nf
have f9 : |1 - (-1)| = |(1 - L) + (L - (-1))| := by
        ring_nf
have f10 : |(1 - L) + (L - (-1))| \le |(1 - L)| + |(L - L)|
         (-1)) | := by apply abs_add
have f11 : |(L - (-1))| = |-((-1) - L)| := by ring_nf
have f12 : |-((-1) - L)| = |((-1) - L)| := by apply
have f13 : (2 : \mathbb{R}) < 1/2 + 1/2 := by linarith [f8, f9,
           f10, f11, f12, f7, f1, f2]
norm_num at f13
```

3 Step 3: Natural Language Proof

Once more, in natural language.

3.1 Natural Language Proof

Theorem: The sequence 1, -1, 1, -1, ... does not converge. **Proof by contradiction:**

- 1. Suppose the sequence converges to some limit L
- 2. Set $\varepsilon = 1 / 2$ and apply the convergence definition
- 3. This guarantees some N such that |a n L| < 1 / 2 for all $n \ge N$
- 4. Consider two specific indices: n = 2N and n = 2N+1 (both are $\geq N$)
 - a $(2N) = (-1)^{2} (2N) = 1$, so |1 L| < 1 / 2
 - a $(2N+1)=(-1)^{(2N+1)}=-1$, so |-1-L|<1/2.
- 5. But then: 2 = |1 (-1)| = |(1 L) + (L + 1)|
 - $\leq |1 L| + |L + 1|$ [triangle inequality]
 - = |1 L| + |-1 L| [algebraic manipulation]
 - < 1/2 + 1/2 = 1 [from steps above]
- 6. This gives us 2 < 1, which is impossible. QED

3.2 Key Insights

- Geometric intuition: Any proposed limit L must be within 1/2 of both 1 and -1 simultaneously, but these values are distance 2 apart
- Critical choice: We chose ε = 1/2 strategically; larger values like ε = 1.5 wouldn't work since L = 0 could satisfy the constraints.
- Triangle inequality: The key step uses |a + b| ≤|a| + |b| to convert the distance between sequence values into a sum of distances from the limit.

This proof exemplifies how rigorous definitions resolve historical mathematical debates and provide clear criteria for convergence.